

Using and reusing coherence to realize quantum processes

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Coherent superposition is a key feature of quantum mechanics that underlies the advantage of quantum technologies over their classical counterparts. Recently, coherence has been recast as a resource theory in an attempt to identify and quantify it in an operationally well-defined manner. Here we study how the coherence present in a state can be used to implement a quantum channel via incoherent operations and, in turn, to assess its degree of coherence. We introduce the robustness of coherence of a quantum channel—which reduces to the homonymous measure for states when computed on constant-output channels—and prove that: i) it quantifies the minimal rank of a maximally coherent state required to implement the channel; ii) its logarithm quantifies the amortized cost of implementing the channel provided some coherence is recovered at the output; iii) its logarithm also quantifies the zero-error asymptotic cost of implementation of many independent copies of a channel. We also consider the generalized problem of imperfect implementation with arbitrary resource states. Using the robustness of coherence, we find that in general a quantum channel can be implemented without employing a maximally coherent resource state. In fact, we prove that *every* pure coherent state in dimension larger than 2, however weakly so, turns out to be a valuable resource to implement *some* coherent unitary channel. We illustrate our findings for the case of single-qubit unitary channels.

Quantum theory has profoundly altered the way we view the physical world. Since its con-

crete formulation roughly a century ago the theory has produced startling predictions, such as quantum entanglement [20] and the no-cloning theorem [66] to name just two, that bear no analog in our classical theories of physics and information processing. The underlying cause of departure of quantum theory from classical ways of thought—and perhaps its only mystery [25]—is the principle of superposition. Indeed, the fact that quantum systems can exist in a superposition state underlies every advantage furnished by quantum technologies over their classical counterparts, from thermodynamics [46, 47, 23, 53, 44] and non-classicality [43], condensed matter [69, 49], metrology [30, 62, 19] and atomic clocks [48], quantum simulation [14, 28] and computation [26, 45, 35, 3, 52] to communication [32, 31, 36], and has even been employed in the description of fundamental biological processes [22, 42, 15, 41].

The aforementioned advantages render quantum superposition a precious resource; having access to it allows one to perform tasks that are otherwise impossible. Hitherto, the focus has been primarily on the quantification and interconversion of *static coherence*—the degree of superposition present in a state [4, 65, 1]. This overlooks the principal goal of a resource, namely that it can be spent in order to perform non-trivial tasks. Some initial steps in this direction have been undertaken for some particular tasks [4, 5, 54, 7, 29, 35]. Here we establish a general framework for converting static coherence into *dynamic coherence*—the ability to generate coherence itself. Specifically, we quantify the amount of static coherence used to implement an arbitrary quantum process and show that in general some static coherence can be reused after the implementation. We prove that both these

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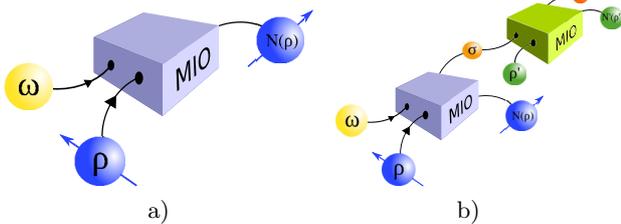


Figure 1: Different protocols for implementing a quantum channel \mathcal{N} using a MIO \mathcal{M} and a coherent resource ω . **a)** The implementation destroys completely the input resource, $\mathcal{M}(\omega \otimes \rho) = \mathcal{N}(\rho)$. **b)** After the implementation, a degraded resource σ is recovered and ready to be recycled: $\mathcal{M}(\omega \otimes \rho) = \sigma \otimes \mathcal{N}(\rho)$. For example it can be used to implement another channel \mathcal{N}' : $\mathcal{M}'(\sigma \otimes \rho') = \tau \otimes \mathcal{N}'(\rho')$.

quantities are quantified by the same measure that we call channel robustness of coherence, in analogy to that of states.

The coherence of a state can be cast within the general framework of a quantum resource theory [8]. Building upon the works of Baumgratz *et al.* [4] and Åberg [1] (see also [11]), the construction of a full-fledged resource theory of coherence, much like those of entanglement [56, 63, 39, 21], asymmetry [33, 51], and athermality [9, 10, 37, 24, 34, 55], has received increased attention in recent years. In the general framework of quantum resource theories one first identifies which states and operations are free and then studies how resource states can be quantified, manipulated, and interconverted among each other. In the resource theory of coherence the set Δ of free states—called *incoherent states*—consists of all the states $\delta \in \mathcal{S}(\mathcal{H})$ that are diagonal in some fixed basis $\{|i\rangle\}_{i=0}^{d-1}$ of \mathcal{H} , i.e. $\Delta = \{\sum_i p_i |i\rangle\langle i| : \sum_i p_i = 1\}$, determined by the problem under investigation. Free operations are those that map the set of free states to itself, hence do not generate coherence. Unlike the resource theories of entanglement, asymmetry and athermality, where the set of free operations is uniquely identified on operational grounds, no such consensus exists within the resource theory of coherence [17]. The largest class of free operations are the *maximally incoherent operations* (MIOs) [1], comprising all completely positive and trace-preserving (CPTP) quantum channels \mathcal{M} such that $\mathcal{M}(\Delta) \subset \Delta$. A subset of MIOs are *incoherent operations* (IOs) [4] which

consist of all quantum channels \mathcal{M} admitting a Kraus representation with operators K_α such that $K_\alpha \Delta K_\alpha^\dagger \subset \Delta$ for all α . Finally, a *coherence measure* [4] is a functional $C : \mathcal{S}(\mathcal{H}) \rightarrow \mathbb{R}_{\geq 0}$, quantifying the amount of coherence present in a state $\rho \in \mathcal{S}(\mathcal{H})$ that satisfies the following two conditions; (i) faithfulness, which means that $C(\delta) = 0$ for all incoherent $\delta \in \Delta$, and (ii) monotonicity, $C(\rho) \geq C(\mathcal{M}(\rho))$, for all free operations \mathcal{M} . Examples of coherence measures include the relative entropy of coherence [4], the robustness of coherence [54], the coherence rank (coherence number) [43] and the ℓ_1 -norm of coherence [4]; the latter two are valid measures only under IOs as free operations, but not MIOs (see Appendix A.1 for further details, and [61] for an extensive review). Finally, the *maximally coherent state* in dimension d , irrespectively of the coherence measure of choice, is the state $|\Psi_d\rangle = \frac{1}{\sqrt{d}} \sum_{j=0}^{d-1} |j\rangle$, which we will henceforth call *cosdit*, reserving the term *cosbit* for the unit of quantum coherence—the maximally coherent qubit.

In this work, we go beyond the resource theory of states and investigate how the coherence present in a state can be used to implement a quantum channel. We thereby assess the coherence of the channel itself, in the spirit of channel resource theory [5, 11]. Previous approaches addressed only the cases where the available operations are IOs and asymptotically many rounds of implementation are allowed [4, 18, 5, 11]. Here we benchmark the ultimate performance of implementation by using MIOs as free operations. Moreover, we mainly consider single-shot channel implementation and discuss its relevance on asymptotics in Sec. 1.3. We find necessary and sufficient conditions for a channel to be implemented with cosdit resource states of any chosen dimension (Fig. 1a), significantly advance the results of previous works on the same subject [4, 18, 5]. The minimal dimension of such a cosdit is used to quantify the *simulation cost* of the channel. In addition, we consider the setting where the input cosdit is recovered in a degraded form at the output, after implementing the required channel. This setting is particularly appealing since the initial resource can be recycled for further use (Fig. 1b) and it naturally gives rise to the notion of an *amortized simulation cost* as the difference between the input and output coherence, which quantifies the resources that are

actually consumed by the implementation.

We begin by introducing the robustness of coherence of quantum channels in Sec. 1.1 and single it out as a natural coherence measure for channels, as it quantifies the simulation and amortized costs of a channel with cosdit resource states. In Sec. 1.2 we consider instead the implementation of channels with arbitrary resource states, providing a semidefinite program to compute the minimum implementation error and the amortized cost. We then introduce a family of non-maximally coherent resource states that allow for the implementation of arbitrary quantum channels and prove that in fact *every* pure coherent state of dimension $d > 2$ is useful for the *exact* implementation of *some* coherent unitary channel. Finally, we focus on the case of qubit unitaries: (i) we prove that, among qubit resource states, a cosbit is necessary and sufficient to implement any coherent qubit unitary; (ii) for the case of qutrit resource states we provide a full characterization with the aid of numerics; (iii) we show that MIOs, unlike IOs [5], allow for coherence recycling. Eventually, in Sec. 1.3 we discuss the relation of our findings to the problem of asymptotic reversibility between coherence cost and coherence generating capacity of quantum channels. For ease of exposition, we defer the proofs of all Theorems, Lemmas and Propositions to Appendices.

1 Results

1.1 MIO implementation of quantum channels using maximally coherent resources

The main goal of the present article is to quantify the resources required to implement (or simulate) an arbitrary quantum channel via MIOs by making use of coherent input states. There are many inequivalent ways to quantify the coherence of a given state. Indeed, a state can be more resourceful than another according to a given measure, while the opposite can happen according to a different one. This is due to the existence of *incomparable* resources, i.e., pairs of states that cannot be interconverted in either direction via MIOs¹.

¹For instance, the cosbit cannot be transformed via MIO into the flagpole state $|\varphi_{\frac{2}{3}}\rangle$ (defined below in Eq. (22)), since the latter has a greater robustness of coherence $C_R(\varphi) = \frac{5}{3} > 1 = C_R(\Psi_2)$; and at the same time

Nevertheless, in this section we wish to establish general results which hold irrespectively of the particular coherence measure of choice: we will do this at the expense of quantifying the required resource in a somewhat coarse fashion, namely by the smallest size k of a cosdit $|\Psi_k\rangle$ that allows to implement the channel. Indeed, cosdits are maximally coherent states irrespectively of how coherence is quantified, since any state of the same (or lower) coherence rank can be obtained from them via MIOs [4]. Restricting our input resources to cosdits might seem a limitation. Nevertheless, we will show that this setting has clear benefits and that it leads to a coherence measure for channels that does take real values.

Our starting point is to introduce the channel robustness:

Definition 1 *The robustness of coherence of a quantum channel \mathcal{N} , $C_R(\mathcal{N})$, is defined by*

$$1 + C_R(\mathcal{N}) := \min \{ \lambda : \mathcal{N} \leq \lambda \mathcal{M}, \mathcal{M} \in \text{MIO} \}, \quad (1)$$

where the inequality $\mathcal{N} \leq \lambda \mathcal{M}$ is understood as completely-positive ordering of super operators, i.e., $\lambda \mathcal{M} - \mathcal{N}$ is a CP map. The smoothed version of this quantity, called ϵ -robustness of coherence, is given by

$$C_R^\epsilon(\mathcal{N}) = \min \left\{ C_R(\mathcal{L}) : \frac{1}{2} \|\mathcal{N} - \mathcal{L}\|_\diamond \leq \epsilon \right\}, \quad (2)$$

where $\|\cdot\|_\diamond$ denotes the diamond norm [2, 64].

Recall that the diamond norm is a well-behaved distance which is furthermore endowed with an operational meaning: it quantifies how well one can physically discriminate between two quantum channels [57]. Also note that definition 1 reduces to the robustness of coherence of states [54] when applied to the constant channel $\mathcal{N}_\sigma(\rho) = \sigma$, i.e., $C_R(\mathcal{N}_\sigma) = C_R(\sigma)$. It is straightforward to see (Appendix A.2) that, just as in the case of states, the robustness of coherence of a channel \mathcal{N} quantifies the minimum amount of noise, in the form of another channel, that we need to add to \mathcal{N} such that the resulting channel is MIO. In Appendix A.2 we also show that this quantity is a proper coherence measure for channels under MIO and that it can be easily formulated

the flagpole $|\varphi_{\frac{2}{3}}\rangle$ cannot be transformed via MIO into the cosbit by Lemma 12.

as a semidefinite program. Such formulation allows for an efficient numerical computation of this measure and, more importantly, thanks to its dual form, it facilitates the proof of the following theorem. In it, $P_R(\mathcal{N})$ [50, 13] and $\widehat{P}_R(\mathcal{N})$ [27, 13] are two alternative, seemingly inequivalent definitions of the cohering power:

$$\begin{aligned} P_R(\mathcal{N}) &= \max_{\rho \in \Delta} C_{LR}(\mathcal{N}(\rho)) \\ &= \max_i \log(1 + C_R(\mathcal{N}(|i\rangle\langle i|))), \end{aligned} \quad (3)$$

$$\widehat{P}_R(\mathcal{N}) = \max_{\rho \in \mathcal{S}(\mathcal{H})} (C_{LR}(\mathcal{N}(\rho)) - C_{LR}(\rho)), \quad (4)$$

with $C_{LR}(\rho) = \log(1 + C_R(\rho))$ being the log-robustness of a quantum state ρ [54] (see Appendix A.1); in particular, for a cosdit it holds $C_{LR}(\Psi_k) = \log k$.

Theorem 2 *The log-robustness of coherence of a channel \mathcal{N} , defined in the following equation, can be expressed as*

$$C_{LR}(\mathcal{N}) := \log(1 + C_R(\mathcal{N})) = P_R(\mathcal{N}) = \widehat{P}_R(\mathcal{N}) \quad (5)$$

and it is additive under tensor product of channels.

Note that throughout the article the base-2 log is employed. Moreover, it can be shown that the (smooth) log-robustness of a channel \mathcal{N} equals the (smooth) maximum relative entropy between \mathcal{N} and a MIO \mathcal{M} , minimized over all \mathcal{M} , that we define in Appendix A.2. We then see that Theorem 2 states that the log-robustness of a channel has an operational meaning as its cohering power. Moreover, the robustness and log-robustness play a prominent role in quantifying the cost and consumption of coherence in the implementation of a channel via MIOs, respectively, as we now explain.

Let us analyze the implementation setting of Fig. 1a, when cosdits are employed as resources:

Definition 3 *A MIO simulation of a quantum channel $\mathcal{N} : A \rightarrow B$ up to error ϵ with a cosdit Ψ_k is a MIO $\mathcal{M} : R \otimes A \rightarrow B$ that satisfies*

$$\frac{1}{2} \|\mathcal{N} - \mathcal{M}(\Psi_k \otimes \cdot)\|_{\diamond} \leq \epsilon. \quad (6)$$

The simulation cost of \mathcal{N} , denoted $C_{\text{sim}}^{\epsilon}(\mathcal{N})$, is the smallest $\log k$ satisfying Eq. (6), i.e., the minimal coherence rank of the resource allowing for a MIO simulation of \mathcal{N} .

Henceforth, $C_{\text{sim}}(\mathcal{N})$ implies $\epsilon = 0$, i.e., exact implementation. We are now equipped to address one of the main objectives of this article and give an exact expression for the simulation cost of a channel in terms of its smoothed robustness of coherence:

Theorem 4 *For any quantum channel \mathcal{N} it holds*

$$C_{\text{sim}}^{\epsilon}(\mathcal{N}) = \log \lceil 1 + C_R^{\epsilon}(\mathcal{N}) \rceil, \quad (7)$$

where $\lceil \cdot \rceil$ is the ceiling function.

Theorem 4 can be seen as one-shot coherence dilution in the resource theory of channels: a maximally coherent resource is completely consumed to generate a target one with smaller coherence-generating capability. Indeed, this includes as a special case the one-shot coherence dilution of states recently studied in [67] and generalizes the criterion found in [17] for transformations of cosdits Ψ_k to pure states $|\phi\rangle = \sum_i \sqrt{p_i} |i\rangle$:

$$\sum_i \sqrt{p_i} \leq \sqrt{k}. \quad (8)$$

Thanks to Theorem 4, we can now determine the minimal coherence rank of a cosdit necessary for a MIO implementation of a channel². This provides a somewhat coarse measure of the implementation cost, as cosdits come in discretized units, the smallest being $k = 2$. In order to quantify the actual resource consumed in the process, in the remaining part of this subsection we investigate to which extent some of the coherence of the input resource can be recovered after the channel implementation. For this purpose let us now focus on the setting of Fig. 1b, where the resource, initially in a state ω , is recovered in a degraded form σ after fueling the MIO implementation of the target channel $\mathcal{N}(\rho)$. We quantify the minimal resource consumed in this process by

²We also note that, since the the cosdit is the maximally coherent state with a given coherence rank, Theorem 4 implies that a generic input resource state ω will not be able to ϵ -simulate a channel if its coherence, as measured by the coherence rank (or by coherence number if the resource is a mixed state, both defined in Appendix A.1), is smaller than the simulation cost: $C_{\text{rank}}(\omega) < \lceil 1 + C_R^{\epsilon}(\mathcal{N}) \rceil$. Indeed, the latter inequality implies that a cosdit ψ'_k of the same rank $k' = C_{\text{rank}}(\omega)$ cannot simulate the channel; moreover, by construction ω is less coherent than ψ'_k , since it can be obtained from the cosdit via $IO \subset MIO$ [4]; hence ω cannot simulate a channel already not simulable with ψ'_k .

the difference between the coherence of ω and σ , when both the input and output states come in standard coherence units, i.e., ω and σ are cosdits.

Definition 5 The ϵ -error amortized cost of a quantum channel \mathcal{N} is given by

$$\begin{aligned} C_{\text{amo}}^\epsilon(\mathcal{N}) &:= \inf (C_{LR}(\Psi_k) - C_{LR}(\Psi_m)) \\ &= \inf \left(\log \frac{k}{m} \right) \\ \text{s.t. } \mathcal{M} &\in \text{MIO}, \\ \mathcal{M}(\Psi_k \otimes \cdot) &= \Psi_m \otimes \mathcal{L}(\cdot), \\ \frac{1}{2} \|\mathcal{N} - \mathcal{L}\|_\diamond &\leq \epsilon, \end{aligned} \quad (9)$$

where we recall that $C_{LR}(\Psi_k) = \log k$. In other words the optimization is over all channels \mathcal{L} that are ϵ -close to the target channel \mathcal{N} and that can be implemented via a MIO $\mathcal{M} : R \otimes A \rightarrow S \otimes B$ with cosdit input and output resource states, respectively Ψ_k and Ψ_m , $k \geq m$.

Note that tensor-product structure at the output of the simulation allows complete freedom in reusing σ and $\mathcal{N}(\rho)$ afterwards; an entangled output would unnaturally constrain the recycling operations. For example, it would not allow the implementation of a sequence of channels applied on-the-fly to the same system, i.e., $\mathcal{N}_n \circ \dots \circ \mathcal{N}_2 \circ \mathcal{N}_1$, whereas this is allowed by Eq. (9), see Corollary 9 below.

Interestingly, the amortized cost defined in Eq. (9) can be related to the log-robustness of a channel. In order to show it, let us first discuss the exact implementation of a channel via MIOs with coherence recycling, which amounts to taking $\epsilon = 0$ above and $\mathcal{L} \equiv \mathcal{N}$. In this case, thanks to Theorem 4, it is possible to estimate the robustness of coherence left in the resource after the implementation:

Theorem 6 For a quantum channel $\mathcal{N} : A \rightarrow B$, there exists a MIO $\mathcal{M} : R \otimes A \rightarrow S \otimes B$ such that $\mathcal{M}(\Psi_k \otimes \cdot) = \sigma \otimes \mathcal{N}(\cdot)$ if and only if

$$C_{LR}(\mathcal{N}) \leq C_{LR}(\Psi_k) - C_{LR}(\sigma). \quad (10)$$

Note that the bound in Eq. (10) can always be attained by some resource state σ_0 that, however, will not be a maximally coherent. If we do impose that the output resource is a cosdit Ψ_m , we obtain the following result:

Corollary 7 Given a quantum channel \mathcal{N} and an integer $k \geq 1 + C_R(\mathcal{N})$, there always exists a MIO implementation of \mathcal{N} that takes a cosdit resource of coherence rank k and returns a degraded resource in the form of a cosdit of rank $m = \lfloor \frac{k}{1 + C_R(\mathcal{N})} \rfloor$, i.e.

$$1 + C_R(\mathcal{N}) \leq \frac{k}{m} \leq (1 + C_R(\mathcal{N})) \left(1 + \frac{1}{m} \right). \quad (11)$$

That is, demanding a cosdit at the output requires an overhead of at most $O(1/m)$ with respect to the optimal ratio that can be attained with a non-maximally coherent output resource; moreover, this overhead can be made arbitrarily small by simply providing a higher-rank cosdit resource at the input. This implies that the amortized cost of a channel is equal to its log-robustness and the same straightforwardly extends to the approximate case, as proved by the following theorem:

Theorem 8 For any quantum channel \mathcal{N} it holds

$$C_{\text{amo}}^\epsilon(\mathcal{N}) = C_{LR}^\epsilon(\mathcal{N}), \quad (12)$$

where

$$C_{LR}^\epsilon(\mathcal{N}) = \min \left\{ C_{LR}(\mathcal{L}) : \frac{1}{2} \|\mathcal{N} - \mathcal{L}\|_\diamond \leq \epsilon \right\}. \quad (13)$$

This second key result, together with Theorem 4, establishes the robustness of coherence of a channel as the correct measure to quantify the cosdit resources necessary to implement the channel using a single MIO. Eq. (7) determines the minimum coherence rank required at the input, while Eq. (12) determines the minimum fraction of input coherence that is actually used in the process. Note however that the latter is a lower bound on the actual coherence consumed when a limited amount of resource can be employed at the input.

Finally, restricting again to the zero-error case, Corollary 7 paves the way for the exact MIO implementation of arbitrary sequences of channels, as depicted in Fig. 1b.

Corollary 9 Any succession of channels $\mathcal{N}_1, \mathcal{N}_2, \dots, \mathcal{N}_n$ can be implemented on-the-fly by concatenation of MIO implementations at an

asymptotically optimal amortized cost

$$\sum_{i=1}^n C_{\text{amo}}(\mathcal{N}_i) \leq \log \frac{k}{k'} \leq \sum_{i=1}^n C_{\text{amo}}(\mathcal{N}_i) + \frac{n}{k'}, \quad (14)$$

where k and $k' \leq k$ are the coherence ranks of the input and output cosbits of the entire protocol.

Note that the additivity of the zero-error log-robustness under tensor product, Theorem 2, already implies $C_{\text{amo}}(\mathcal{N}_1 \otimes \dots \otimes \mathcal{N}_n) = \sum_{i=1}^n C_{\text{amo}}(\mathcal{N}_i)$. However, Corollary 9 is more general and includes this result as a special case. Indeed, it allows the exact implementation of *arbitrary sequences* of channels, and not just their tensor product, at an amortized cost equal to the sum of the single amortized costs of each channel. In particular, this allows the implementation of a concatenation of channels on-the-fly and it justifies the choice of a tensor-product structure at the output of the recycling process. Finally, from Theorem 2 we know that the log-robustness of coherence also quantifies the cohering power of a quantum channel. Hence we conclude that the exact amortized cost of a channel coincides with its cohering power.

Example: Qubit unitaries

Let us now focus on implementing the simplest kind of channels, qubit unitary gates. A unique representation of a qubit unitary up to incoherent symmetries is given by [5]

$$U_\theta = \begin{pmatrix} c & -s \\ s & c \end{pmatrix}, \quad (15)$$

where $c = \cos \theta$, $s = \sin \theta$ and $0 < \theta \leq \frac{\pi}{4}$. In this case we can compute all the quantities defined above:

$$C_{LR}(U_\theta) = \log(1 + \sin 2\theta) = C_{\text{amo}}(U_\theta), \quad (16)$$

$$C_{\text{sim}}(U_\theta) = \begin{cases} 0 & \theta = 0, \\ 1 & \text{otherwise.} \end{cases}, \quad (17)$$

as can be readily checked by noticing that the robustness of a qubit state is equal to its ℓ_1 -norm of coherence (see Eq. (34)): $C_R(U_\theta|i\rangle\langle i|U_\theta^\dagger) = C_{\ell_1}(U_\theta|i\rangle\langle i|U_\theta^\dagger) = 2cs = \sin 2\theta$.

Eventually, in contrast with the case of IOs [5], where the implementation of qubit unitaries consumes the entire cosbit resource, Theorem 6 ensures that MIOs do allow for coherence recycling.

More precisely, there exists a MIO \mathcal{M} such that $\mathcal{M}(\Psi_2 \otimes \rho) = \sigma \otimes U_\theta \rho U_\theta^\dagger$ for any qubit state ρ if and only if the output resource state σ has coherence

$$C_R(\sigma) \leq \frac{1 - \sin 2\theta}{1 + \sin 2\theta}. \quad (18)$$

1.2 Arbitrary resources for MIO implementation of channels

Let us now go beyond the assumptions of Sec. 1.1 by considering a scenario where non-maximally coherent states are employed as resources (see [59] for an analogous result in entanglement theory). In particular, we want to study under which conditions the MIO implementation of a quantum channel in the settings of Fig. 1a and Fig. 1b is still possible.

We begin by introducing a semidefinite program to assess the performance of an arbitrary resource at implementing any target channel; this program yields the best approximate MIO implementation of the target channel with a given resource state (Fig. 1a), as measured by the diamond norm.

Proposition 10 *The smallest diamond norm error for the implementation of a quantum channel $\mathcal{N} : A \rightarrow B$ via a MIO $\mathcal{M} : R \otimes A \rightarrow B$ with a coherent resource $\omega \in R$ is given by the following semidefinite program:*

$$\begin{aligned} \min \quad & \lambda \\ \text{s.t.} \quad & J_{\mathcal{M}} \text{ is the Choi matrix of } \mathcal{M} \in \text{MIO} \\ & J_{\mathcal{E}} = \text{tr}_R \left(\left(\omega^t \otimes \mathbb{1}_A \otimes \mathbb{1}_B \right) J_{\mathcal{M}} \right) \\ & Z \geq 0 \\ & Z \geq J_{\mathcal{N}} - J_{\mathcal{E}} \\ & \lambda \mathbb{1}_A \geq \text{tr}_B Z, \end{aligned} \quad (19)$$

where $J_{\mathcal{N}}$ is the Choi matrix of \mathcal{N} , ω^t denotes the transpose of ω , and $J_{\mathcal{E}}$ that of its implementation, $\mathcal{E} = \text{tr}_B \mathcal{M}(\omega \otimes \cdot)$. Recall that the Choi matrix $J_{\mathcal{M}}$ of a MIO \mathcal{M} as in Eq. (19), is fully characterized by $J_{\mathcal{M}} \geq 0$, $\text{tr}_B J_{\mathcal{M}} = \mathbb{1}_{RA}$ and $\text{tr}(|i\rangle\langle i|_{RA} \otimes |j\rangle\langle k|_B J_{\mathcal{M}}) = 0 \forall i$ and $\forall j \neq k$.

In Appendix A.4 we provide a simpler semidefinite program for the case of unitary channels, where the precision of the simulation is assessed in terms of the average gate fidelity $f(U, \mathcal{N})$ [38], rather than by the diamond distance. Specifically, we compute the entanglement fidelity

$F(U, \mathcal{N}) = \text{tr } J_U J_{\mathcal{N}}$ (where J_U and $J_{\mathcal{N}}$ are the Choi matrices of U and \mathcal{N} , respectively), which fulfills $F(U, \mathcal{N}) = \frac{(d+1)f(U, \mathcal{N}) - 1}{d}$, where d is the dimension of the Hilbert space on which the channels act [38]. From now on we will refer to it as ‘‘gate fidelity’’.

Regarding instead the recycling setting of Fig. 1b, thanks to Theorem 8 it is straightforward to write a semidefinite program for the ϵ -error amortized cost of a channel with cosdit input resource, as shown in Fig. 5. If, in addition, we ask for the maximum robustness of coherence left in the output resource when a non-maximally coherent input resource is employed, we end up with the following optimization problem:

Proposition 11 *The maximum coherence left in the resource $\sigma \in S$ after the implementation of a quantum channel $\mathcal{N} : A \rightarrow B$ via a MIO $\mathcal{M} : R \otimes A \rightarrow S \otimes B$ and a coherent resource $\omega \in R$ up to error ϵ in diamond norm is:*

$$\begin{aligned} \max C_R(\sigma) \\ \text{s.t. } \mathcal{M}(\omega \otimes \cdot) = \sigma \otimes \mathcal{L}(\cdot) \\ \frac{1}{2} \|\mathcal{N} - \mathcal{L}\|_{\diamond} \leq \epsilon. \end{aligned} \quad (20)$$

This problem captures exactly the setting of Fig. 1b with arbitrary input resource. Note, however, that it cannot be formulated as a semidefinite program, since the tensor-product constraint is non-linear in the optimization variables σ and \mathcal{L} . One can devise alternative semidefinite programs by relaxing the constraints in Eq. (20), as we discuss in Appendix A.5.

The semidefinite program in Proposition 10 and that for the amortized cost allow for a thorough numerical analysis of our problem, see Figs 3-5. Nevertheless, before analysing the numerical results for qubits, let us present some general, analytical results on MIO implementation of channels with non-maximally coherent resources. As a start, according to Theorem 4, any resource state that is MIO-convertible into a cosdit of coherence rank d can be trivially used to implement any quantum channel of log-robustness smaller than $\log d$. The following lemma provides us with a simple necessary condition for the existence of pure state transformation via MIO.

Lemma 12 *Assume that a pure state ψ is transformed to a pure state ϕ via MIO, then*

$$\lambda_1(\psi) \leq \lambda_1(\phi), \quad (21)$$

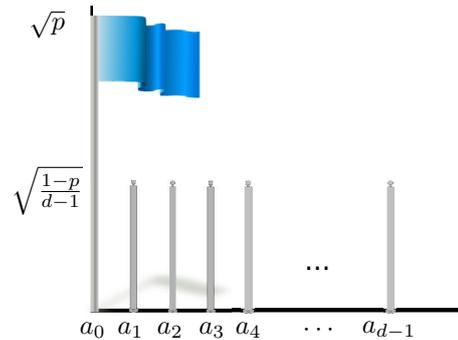


Figure 2: Schematic depiction of the coefficients $\{a_j\}_{j=0}^{d-1}$ of a flagpole state $|\varphi_p\rangle$ in the incoherent basis.

where $\lambda_1(\rho) = \max_i \langle i|\rho|i\rangle$ is the largest diagonal entry of an operator ρ , which coincides with its fidelity of coherence [60, 58] on pure states.

Condition (21) is also sufficient if the target state is maximally coherent, since if $\lambda_1(\psi) \leq \frac{1}{d}$ then there exists a IOCMIO that transforms ψ into Ψ_d due to the majorization criterion [65].

Hence, for any channel \mathcal{N} with $C_{LR}(\mathcal{N}) \leq \log d$ there exists a continuous family of resource states of dimension d that allow for the exact MIO simulation of it via conversion to cosdits. We now ask whether there exist any channels that can be implemented using pure resource states $|\omega\rangle$ that are not convertible into cosdits, i.e., $\lambda_1(\omega) > 1/d$. To this end let us define a special class of states that will be useful in this context.

Definition 13 *A d -dimensional flagpole state is a pure state of the form*

$$|\varphi_p\rangle = \sqrt{p}|0\rangle + \sqrt{\frac{1-p}{d-1}} \sum_{j=1}^{d-1} |j\rangle, \quad \text{with } \frac{1}{d} \leq p \leq 1. \quad (22)$$

The structure of flagpole states, shown in Fig. 2, endows them with several useful properties. Thanks to the majorization criterion [65], namely that the pure state transformation $\phi \xrightarrow{\text{IO}} \psi$ is possible if $\mathcal{D}(\psi) \succ \mathcal{D}(\phi)$ (where \mathcal{D} is the complete dephasing map), it is easy to show that: i) for all pure states ϕ we can transform $\phi \mapsto \varphi_p$ with a specific value of $p \geq F_C(\phi)$ (see Eq. (65)); ii) conversely, we can transform $\varphi_p \mapsto \phi$ for all ϕ such that $p \leq F_C(\phi)$. In other words, φ_p is the most coherent state of fixed coherence rank with fidelity of coherence larger than or equal to p . Moreover, as explained above, a d -dimensional flagpole state with $p \leq \frac{1}{d-1}$ can be converted

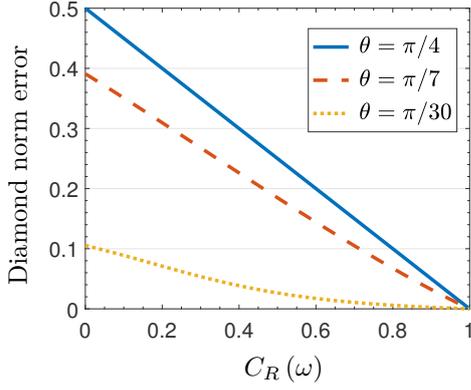


Figure 3: Minimum diamond norm distance between the qubit unitary U_θ and a MIO implementation of it with pure qubit resource state $|\omega\rangle$ vs. robustness of coherence of the latter. For $\theta = \pi/30$ (dotted orange line), $\theta = \pi/7$ (dashed red line) and $\theta = \pi/4$ (solid blue line). Note that an exact implementation is possible only with a cosbit resource state.

into a $(d-1)$ -dimensional cosdit and thus trivially implements any channel of log-robustness smaller or equal than $\log(d-1)$. For all other flagpoles with $p > \frac{1}{d-1}$, which cannot be converted to a cosdit of rank $d-1$, the following theorem holds.

Theorem 14 *For any quantum channel $\mathcal{N} : A \rightarrow B$, if*

$$p \leq \frac{1}{1 + C_R(\mathcal{N})}, \quad (23)$$

then there exists a MIO $\mathcal{M} : R \otimes A \rightarrow B$ such that $\mathcal{M}(\varphi_p \otimes \rho) = \mathcal{N}(\rho)$ for all states ρ , where φ_p is a d -dimensional flagpole state, $d > |B|$.

This proves that any pure resource state in dimension larger than 2 is useful for the exact implementation of some coherent unitary channel via MIO. Indeed, any such state can be converted to a flagpole, which in turn can be used for the implementation of a coherent channel, in particular a unitary. Note also that in general the bound on p in Theorem 14 does not single out all flagpoles that can implement a given channel, since it relies on a specific ansatz (see Appendix A.6 for details).

Example: Qubit unitaries

Analogously to the previous section, we now address the implementation of qubit unitaries when non-maximally coherent resources are available.

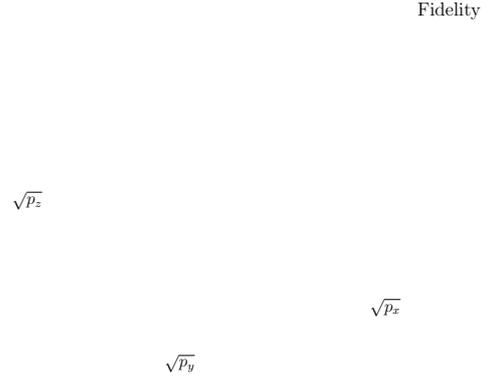


Figure 4: Gate fidelity for the MIO implementation of a qubit unitary U_θ with $\theta = \pi/14$ using a generic qutrit ($d = 3$) resource $|\psi\rangle = \sqrt{p_x}|0\rangle + \sqrt{p_y}|1\rangle + \sqrt{p_z}|2\rangle$. The central cross (\times) represents the maximally coherent state. The dashed red line around it encloses states that can be transformed with MIO to cosbits, the blue-solid line encloses the states that allow for an exact implementation ($F = 1$), while the dashed-dotted green line encloses all states that cannot be obtained via MIO by diluting a cosbit. The dotted yellow lines represent the family of flagpole states.

The first question we want to raise is whether it is possible to use a non-maximally coherent pure qubit resource in order to implement a qubit unitary, even one that generates little coherence. In [5] it was proven that, if the free operations are IOs, this is only possible with a cosbit, no matter how coherent the unitary is. This is still the case under MIOs, even for higher-dimensional unitaries, when restricting to qubit resource states:

Proposition 15 *The only pure qubit resource state $|\omega\rangle \in \mathbb{C}^2$ that permits the MIO implementation of some unitary gate of arbitrary dimension is the cosbit.*

As an illustration of this fact, Fig. 3 shows that only a cosbit resource allows for an exact implementation of the qubit unitary U_θ , for several values of θ . Input coherence can be saved only at the expense of allowing for an approximate implementation of the gate. Moreover, in Fig. 4 we give a full characterization of the optimal performance of a MIO simulation, as measured by the gate fidelity, for general qutrit resource states. Without loss of generality we can focus on the upper region of the plot, $p_x, p_y \leq p_z = \lambda_1(\psi)$. Qubit resources are found in the planes defined by $p_x = p_y = 0$, where perfect simulation ($F = 1$) is only reached for cosbits, $p_z = 1/2$. The red dashed line delimits the qutrit states that can be distilled into a

cosbit, $p_z \leq 1/2$, and hence also attain $F = 1$. However, perfect simulation can also be attained with other qutrit states: those that fall below the solid blue line in Fig. 4. In particular, the qutrit state with the highest value of p_z , i.e. the least coherent qutrit as measured by the fidelity of coherence, that allows for perfect simulation is a flagpole. Indeed, this agrees with the predictions of Theorem 14:

Proposition 16 *A d -dimensional flagpole resource state $|\varphi_p\rangle$ allows for the MIO implementation of a qubit unitary U_θ if*

$$p \leq \frac{1}{1 + \sin 2\theta}, \quad (24)$$

where $d \geq 3$ and $0 < \theta \leq \frac{\pi}{4}$. Furthermore, for a qutrit flagpole state $|\varphi_p\rangle \in \mathbb{C}^3$ this is also the highest allowed value of p .

In Fig. 4 all flagpole states are identified by a yellow dotted line that interpolates between the incoherent state $|0\rangle$ and the cosbit. Finally, the set of states that cannot be obtained via MIO from a cosbit, are enclosed by the dashed-dotted green line which is determined by the intersection of the sphere of qutrit states (in the positive octant) and the plane $\langle \Psi_3 | \psi \rangle = 1$, as shown in Eq. (8). Extensive numerical evidence shows that the blue solid line that delimites the region of states that enable an exact MIO implementation is also given by the intersection of the qutrit sphere with a plane of the form $\langle \Phi_\theta | \psi \rangle \leq f(\theta)$, where the constant $f(\theta)$ and normal vector $|\Phi_\theta\rangle$ can be analytically found by imposing that the plane includes the cosbit and the flagpole saturating Eq. (24).

Regarding the case of coherence recycling after the approximate implementation of a qubit unitary, Fig. 5 shows the amortized cost of qubit unitaries for several error thresholds. As expected, the amortized cost increases with the coherence of the target unitary, since less coherence can be recycled when implementing more coherent unitaries. Moreover, if we allow a larger error threshold to implement the same unitary, it is possible to obtain more coherence back for the same input resource at the cost of a worse approximation. When the amortized cost becomes zero, Theorem 8 implies that there exists a MIO that implements the target unitary without consuming any coherence. This happens if and only if the simulation cost of the unitary for the same error threshold is zero.

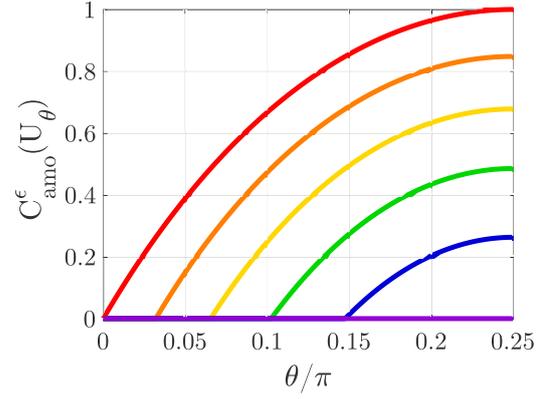


Figure 5: Plot of the ϵ -error amortized cost of the qubit unitary U_θ as vs. θ/π , for several error thresholds $\epsilon \in \{0, 0.1, 0.2, 0.3, 0.4, 0.5\}$ (from top to bottom). The amortized cost is higher for more coherent unitaries but it decreases if larger errors are allowed. In particular, when the amortized cost becomes zero it means that the given unitary can be simulated up to the given error with a MIO, which is not necessarily the identity channel, without using any coherence of the input resource state.

1.3 Comparison with coherence distillation and the reversibility problem

In this work, we have focused on the cost of implementing a channel, but in [5] this was already considered under IOs in conjunction with the capacity to create pure coherence from the given resource channel. This is an important point, because it touches upon the issue of (asymptotic) reversibility of the resource theory. Namely, since in the theory of coherence there exists a unit resource, the cosbit, it is possible to compare the minimum rate of cosbits needed to implement the channel, and the maximum rate of cosbits extractable from the given channel: the theory is asymptotically reversible if and only if these two numbers coincide for all channels. For the latter problem, we find in Ref. [5] a developed theory, even if it was worked out with IOs as free operations. However, the changes required to the definitions, protocols and bounds in [5, Sec. II] (especially to its Thm. 1) are minimal and readily result in the following theorem. Before stating it, let us first recall the generic protocol for coherence generation with many uses of a channel $\mathcal{N} : A \rightarrow B$: i) an initial incoherent state $\rho_0 \in AC$ undergoes the action of the channel, then is post-processed by a MIO $M_1 : BC \rightarrow AC_1$, obtaining the state $\rho_1 = M_1 \circ (\mathcal{N} \otimes \text{id})(\rho_0)$; ii) the

same process is iterated $n - 1$ times, producing the intermediate states $\rho_j = M_j \circ (N \otimes \text{id}) (\rho_{j-1})$, $j = 1, \dots, n$; iii) the iterations stop when the final state ρ_n is ϵ -close to the cosdit $\Psi_2^{\otimes nR}$, i.e., $\text{tr}(\rho_n \Psi_2^{\otimes nR}) \geq 1 - \epsilon$, with R being the resulting coherence generating rate. Then the coherence-generating capacity of \mathcal{N} is defined as follows:

$$C_{\text{gen},\text{MIO}}^{(\infty)}(\mathcal{N}) := \sup_{\epsilon > 0} \limsup_{n \rightarrow \infty} R. \quad (25)$$

Theorem 17 *The asymptotic coherence-generating capacity of a channel $\mathcal{N} : A \rightarrow B$ is given by*

$$C_{\text{gen},\text{MIO}}^{(\infty)}(\mathcal{N}) = \sup_{\rho_{AC}} C_r((\mathcal{N} \otimes \text{id}_C)(\rho_{AC})) - C_r(\rho_{AC}) \quad (26)$$

$$= \sup_C \widehat{P}_r(\mathcal{N} \otimes \text{id}_C) = \mathbb{P}_r(\mathcal{N}), \quad (27)$$

where the suprema are over all auxiliary systems C and the first additionally over states ρ_{AC} on $A \otimes C$. The second formula features the relative entropy cohering power [27, 5] $\widehat{P}_r(\mathcal{N}) := \max_{\rho} C_r(\mathcal{N}(\rho)) - C_r(\rho)$; $\mathbb{P}_r(\mathcal{N}) := \sup_k \widehat{P}_r(\mathcal{N} \otimes \text{id}_k)$ was introduced in [5] as complete relative entropy cohering power. ■

The coherence-generating capacity, despite its nice information theoretic formula, is by no means an easy quantity to compute. Notwithstanding the explicit expression for the relative entropy of coherence, the maximization over the state ρ_{AC} is not necessarily well-behaved. Also, we do not know how large an auxiliary system C is required, or indeed if any at all.

On the other hand, the asymptotic cost of implementing the channel can be expressed using our above results, in particular Theorem 4. Namely, it comes down to

$$C_{\text{sim},\text{MIO}}^{(\infty)}(\mathcal{N}) = \sup_{\epsilon > 0} \limsup_{n \rightarrow \infty} \frac{1}{n} \log(1 + C_R^{\epsilon}(\mathcal{N}^{\otimes n})). \quad (28)$$

There is one special case in which we know what this limit is, namely when for all ρ , $\mathcal{N}(\rho) = \sigma$, i.e. \mathcal{N} is a constant channel. Then, $C_R(\mathcal{N}) = C_R(\sigma)$, and the right hand side of Eq. (28) converges [8] to the relative entropy of coherence, $C_r(\sigma)$. More generally, for a cq-channel

$\mathcal{N}(|i\rangle\langle j|) = \delta_{ij}\sigma_i$, written in the incoherent basis of the input state, the same reasoning yields

$$C_{\text{sim},\text{MIO}}^{(\infty)}(\mathcal{N}) = C_{\text{gen},\text{MIO}}^{(\infty)}(\mathcal{N}) = \max_i C_r(\mathcal{N}(|i\rangle\langle i|)),$$

i.e., the theory is asymptotically reversible when restricted to cq-channels. It is not known under which conditions this happens for general channel resource theories, since the results of [8] cannot be directly applied to them.

2 Discussion

In this work we studied the conversion of static coherence into dynamic coherence under the most general framework, which employs MIOs as the free operations. To that end, we have introduced the robustness of coherence of channels, which includes as a special case the corresponding measure for states, showing its operational meaning in two respects: i) the log-robustness equals the cohering power of the channel; ii) when maximally coherent states are employed as resources, the robustness determines the implementation cost of the channel and the log-robustness characterizes the amortized implementation cost with recycling of coherence at the output, as well as the asymptotic cost of realizing exactly many independent instances of the channel.

We have then considered the more general case of implementation with arbitrary resource states and provided a semidefinite program to find the best approximation to any given channel with any resource. In the same setting, we have analytically demonstrated that any state in dimension larger than 2, however weakly so, is useful for the exact implementation of some coherent unitary channel, owing to the fact that every such state can be converted into a flagpole state and that, in turn, every flagpole is able to realize some channel.

Throughout the article, we have considered the specific case of qubit unitary implementation and studied it in detail, showing a direct application of our general results. Finally, we have related our findings to the problem of asymptotic reversibility, by introducing the asymptotic coherence generating capacity and the asymptotic cost of implementation of a given channel under MIOs. Since both quantities are not easily computable, it re-

mains to be known whether they coincide, i.e. whether the theory is asymptotically reversible.

Our findings significantly advance the state of the art on the theory of coherence of channels, providing an operational measure for its quantification and a good starting point for future research on the subject. In particular, some open questions immediately follow from our results: i) whether a closed expression for the asymptotic implementation cost of channels can be obtained from the smoothed robustness of coherence; ii) whether the coherence left after the approximate implementation of a channel with an arbitrary input resource can be formulated as an efficient optimization problem; iii) whether the set of pure resource states that allow for the exact MIO implementation of an arbitrary unitary is always determined by a linear function of their coefficients in the incoherent basis.

Finally, our results may also spur investigation on other lines of research such as, e.g., exploring similar approaches in the resource theories of non-classicality and athermality, analysing the connection of our findings with those in entanglement theory [6, 21], studying the applicability of coherence theory in an experimental setting and extending further the general framework of channel resource theories with the help of quantum combs [16].

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A Appendix

A.1 Resource theory of coherence: main definitions

In this appendix we provide a brief review of the resource theory of coherence focusing mainly on those elements of the theory that are relevant to our work. For a more exhaustive review of the resource theory of coherence we direct the reader to [61, 40].

The set of *free states*, Δ , in the resource theory of coherence, also referred to as *incoherent* states, form a convex subset of the set of all density matrices, $\mathcal{S}(\mathcal{H})$, and are defined as those density matrices that are diagonal with respect to a chosen orthonormal basis $\{|j\rangle : j = 0, \dots, d-1\}$ of the Hilbert space \mathcal{H} , $\Delta = \left\{ \delta : \delta = \sum_{j=0}^{d-1} \delta_j |j\rangle\langle j| \right\}$. Free operations are those that map the set of free states onto itself; the largest class of free operations conceivable in a resource theory of coherence, and the primary focus of our work, are the so-called *maximally incoherent operations (MIOs)* [1], and they are simply defined as all CPTP operations, \mathcal{M} , that map Δ onto itself. In particular, the Choi-Jamiołkowski matrix, $J_{\mathcal{M}}$, of a MIO operation $\mathcal{M} : A \rightarrow B$ is characterized by the following conditions, in addition to the standard ones:

$$\text{tr}(|i\rangle\langle i| \otimes |j\rangle\langle j'|) J_{\mathcal{M}} = 0, \quad \forall i \forall j \neq j', \quad (29)$$

which is equivalent to requiring that \mathcal{M} does not generate coherence from incoherent states. A well studied subset of MIO are the *incoherent operations (IOs)*, which are defined as those CPTP operations admitting a Kraus decomposition $\mathcal{M}(\rho) = \sum_{\alpha} K_{\alpha} \rho K_{\alpha}^{\dagger}$, where each Kraus operator preserves incoherence, i.e., $K_{\alpha} \Delta K_{\alpha}^{\dagger} \subset \Delta$.

Any state that possesses coherence with respect to the chosen orthonormal basis qualifies as a resource. The quantification of the resourcefulness of any given quantum state is accomplished by a suitable *coherence measure*, $C : \mathcal{S}(\mathcal{H}) \rightarrow \mathbb{R}_{\geq 0}$. A good coherence measure should satisfy the following two conditions: (i) $C(\rho) = 0$ for all $\rho \in \Delta$, and (ii) $C(\rho) \geq C(\mathcal{T}(\rho))$ for all incoherent CPTP maps \mathcal{T} . Another convenient— but not necessary—condition is convexity, $\sum_i p_i C(\rho_i) \geq C(\sum_i p_i \rho_i)$. The following are examples of coherence measures fulfilling all the previous conditions under MIO.

1. *Robustness of coherence* [54]: quantifies the minimum amount of another state ρ' required such that its convex sum with ρ is an incoherent state. It can be cast as a semidefinite program in primal standard form

$$1 + C_R(\rho) = \min\{\lambda : \rho \leq \lambda\sigma, \sigma \in \Delta\}. \quad (30)$$

Its equivalent dual form is given by

$$1 + C_R(\rho) = \max\{\text{tr} \rho S : S \geq 0, S_{jj} = 1 \forall j\}, \quad (31)$$

which holds because strong duality is fulfilled. The robustness is multiplicative under tensor product of states [68]:

$$1 + C_R(\rho_1 \otimes \rho_2) = (1 + C_R(\rho_1))(1 + C_R(\rho_2)).$$

For the sake of completeness, we provide a self-contained proof of this fact below, see Lemma 18. It is convenient to define a logged version of the robustness of coherence as:

$$C_{LR} = \log(1 + C_R(\rho)), \quad (32)$$

which now becomes an additive quantity under tensor product.

2. *Relative entropy of coherence* [4]:

$$C_r(\rho) = S(\mathcal{D}(\rho)) - S(\rho), \quad (33)$$

where $S(\rho) = -\text{tr}(\rho \log \rho)$ is the von Neumann entropy, and $\mathcal{D} : \mathcal{S}(\mathcal{H}) \rightarrow \mathcal{S}(\mathcal{H})$ is the map that erases all the off-diagonal elements of $\rho \in \mathcal{S}(\mathcal{H})$.

3. The ℓ_1 -norm of coherence [4],

$$C_{\ell_1}(\rho) = \sum_{i \neq j} |\rho_{ij}|, \quad (34)$$

is a coherence measure under IOs, but not MIOs [12]. It is popular because it is simple to compute. More importantly, the ℓ_1 -norm tightly bounds the robustness of coherence [54]

$$\frac{C_{\ell_1}(\rho)}{d-1} \leq C_R(\rho) \leq C_{\ell_1}(\rho), \quad (35)$$

where the upper bound becomes an equality for qubits and any pure state. We also note that $\log(1 + C_{\ell_1}(\rho))$ is additive under tensor product [68].

4. *Coherence rank and coherence number* [43]: For pure states the coherence rank is defined as

$$C_{\text{rank}}(\psi) := \min \left\{ r : |\psi\rangle = \sum_{i=1}^r c_i |c_i\rangle \right\}, \quad (36)$$

where $|c_i\rangle$ is an element of the incoherent basis. That is, the coherence rank of $|\psi\rangle$ is the number of non-zero terms that appear when writing $|\psi\rangle$ in the incoherent basis. This coherence measure can be extended by a convex roof construction, to the *coherence number*:

$$C_{\text{rank}}(\rho) := \min_{\rho = \sum_{\alpha} p_{\alpha} \psi_{\alpha}} \max_{\alpha} C_{\text{rank}}(\psi_{\alpha}). \quad (37)$$

We stress that, like C_{ℓ_1} , C_{rank} is a monotone only under IOs, not under MIOs.

Lemma 18 (Zhu et al. [68]) *For any tensor-product state on a bipartite system, it holds*

$$1 + C_R(\rho_1 \otimes \rho_2) = (1 + C_R(\rho_1))(1 + C_R(\rho_2)). \quad (38)$$

Proof. Consider the two values $\lambda_i = 1 + C_R(\rho_i)$ for $i = 1, 2$. Then by Eq. (30) there exist $\sigma_i \in \Delta$ such that $\rho_i \leq \lambda_i \sigma_i$, which implies $\rho_1 \otimes \rho_2 \leq \lambda_1 \lambda_2 \sigma_1 \otimes \sigma_2$, i.e., that $1 + C_R(\rho_1 \otimes \rho_2) \leq \lambda_1 \lambda_2$.

On the other hand, take S_i that optimize Eq. (31) for both local states, i.e., $1 + C_R(\rho_i) = \mathbf{tr}(\rho_i S_i)$ for $i = 1, 2$, and take as an ansatz for the optimizer of the total robustness $S = S_1 \otimes S_2$. It follows that $1 + C_R(\rho_1 \otimes \rho_2) \geq \mathbf{tr}(\rho_1 S_1) \mathbf{tr}(\rho_2 S_2)$. ■

A.2 The robustness of coherence of a channel

In this subsection we prove several properties of the robustness of coherence for quantum channels, introduced in Eq. (1). Let us start by showing that it is a coherence measure for a coherence theory whose free resources are MIOs.

1. *Faithfulness*: $C_R(\mathcal{N}) = 0$ if and only if \mathcal{N} is MIO. This follows straightforwardly from the definition, since one can take $\lambda = 1$ if and only if $\mathcal{M} \equiv \mathcal{N}$.
2. *Monotonicity*: The robustness of coherence of the channel \mathcal{N} monotonically decreases if we pre- and post-process the input and output using MIO channels \mathcal{L}' and \mathcal{L} respectively, i.e., $C_R(\mathcal{L} \circ \mathcal{N} \circ \mathcal{L}') \leq C_R(\mathcal{N})$, $\forall \mathcal{L}, \mathcal{L}'$ is MIO. Indeed, take $\lambda \geq 0, \mathcal{M}$ MIO such that $\lambda \mathcal{M} - \mathcal{N} \geq 0$ and concatenate it with the channels $\mathcal{L}, \mathcal{L}'$. Defining $\mathcal{M}' = \mathcal{L} \circ \mathcal{M} \circ \mathcal{L}'$ MIO it follows that $\mathcal{L} \circ \mathcal{N} \circ \mathcal{L}' \leq \lambda \mathcal{M}'$ as well, so that any λ feasible for \mathcal{N} is also feasible for $\mathcal{L} \circ \mathcal{N} \circ \mathcal{L}'$.
3. *Convexity*: $\sum_i p_i C_R(\mathcal{N}_i) \geq C_R(\sum_i p_i \mathcal{N}_i)$ for any probability distribution $\{p_i\}$ and $\{\mathcal{N}_i\}$ a collection of quantum channels. Indeed for each i take the minimum λ_i such that $\mathcal{N}_i \leq \lambda_i \mathcal{M}_i$, \mathcal{M}_i MIO. Then define $\bar{\lambda} = \sum_i p_i \lambda_i$ and note that $\{\tilde{p}_i = p_i \lambda_i / \bar{\lambda}\}$ is still a probability distribution. By averaging the inequalities over $\{p_i\}$ and rescaling by $\bar{\lambda}$ we get $\sum_i p_i \mathcal{N}_i \leq \bar{\lambda} \sum_i \tilde{p}_i \mathcal{M}_i$, where the latter is still a MIO. We conclude that $1 + C_R(\sum_i p_i \mathcal{N}_i) \leq \bar{\lambda} = \sum_i \lambda_i p_i$, and thus $\sum_i p_i C_R(\mathcal{N}_i) \geq C_R(\sum_i p_i \mathcal{N}_i)$.

Properties 1-3 above straightforwardly extend to the smooth-robustness of Eq. (2) (the faithfulness condition holds up to error ϵ).

Let us now focus on the SDP formulation of the robustness.

Proof of Theorem 2. We want to calculate the dual of the SDP

$$1 + C_R(\mathcal{N}) = \min\{\lambda : \mathcal{M}' \text{ is MIO and } \mathcal{N} \leq \lambda \mathcal{M}' =: \mathcal{M}\},$$

which is equivalent to

$$\begin{aligned} \min \lambda & \\ \text{s.t. } J_{\mathcal{N}} &\leq J_{\mathcal{M}} \\ \text{tr}_B J_{\mathcal{M}} &= \lambda \mathbb{1}_A \\ \text{tr } J_{\mathcal{M}} (|i\rangle\langle i| \otimes |j\rangle\langle k|) &= 0, \quad \forall i, \forall j \neq k, \end{aligned} \tag{39}$$

where $J_{\mathcal{N}}, J_{\mathcal{M}}$ are the Choi matrices of \mathcal{N} and \mathcal{M} , respectively. This means that we need to minimize the following Lagrangian:

$$\begin{aligned} L &= \lambda + \text{tr } X (\text{tr}_B J_{\mathcal{M}} - \lambda \mathbb{1}_A) - \text{tr } Y (J_{\mathcal{M}} - J_{\mathcal{N}}) \\ &\quad + \sum_{i,j \neq k} z_{jk}^i \text{tr } J_{\mathcal{M}} (|i\rangle\langle i| \otimes |j\rangle\langle k|) \\ &= \lambda + \text{tr} (X \otimes \mathbb{1}) J_{\mathcal{M}} - \lambda \text{tr } X - \text{tr } Y (J_{\mathcal{M}} - J_{\mathcal{N}}) \\ &\quad + \text{tr } J_{\mathcal{M}} \left(\sum_{i,j \neq k} |i\rangle\langle i| \otimes Z_i \right) \\ &= \lambda + \text{tr} (X \otimes \mathbb{1}) J_{\mathcal{M}} - \lambda \text{tr } X - \text{tr } Y (J_{\mathcal{M}} - J_{\mathcal{N}}) \\ &\quad + \text{tr} (J_{\mathcal{M}} Z) \\ &= \lambda (1 - \text{tr } X) + \text{tr } J_{\mathcal{M}} (X \otimes \mathbb{1} - Y + Z) + \text{tr} (J_{\mathcal{N}} Y), \end{aligned} \tag{40}$$

where $Z = \sum_i |i\rangle\langle i| \otimes Z_i$ and Z_i is a zero-diagonal matrix, while $Y \geq 0$. In order to avoid that $\min_{\lambda, J_{\mathcal{M}}} L = -\infty$, we need to impose that $\text{tr } X = 1$ (so X is a state) and $Y = X \otimes \mathbb{1} + Z$.

To calculate the dual, we maximize the terms of the Lagrangian which do not contain the variables of the primal form, λ and $J_{\mathcal{M}}$:

$$\max\{\text{tr } J_{\mathcal{N}} Y : \text{tr } X = 1, Y = X \otimes \mathbb{1} + Z \geq 0\}. \tag{41}$$

The objective function is equivalent to

$$\begin{aligned} \text{tr } J_{\mathcal{N}} Y &= \text{tr } J_{\mathcal{N}} (X \otimes \mathbb{1}) + \text{tr} (J_{\mathcal{N}} Z) \\ &= 1 + \text{tr} (J_{\mathcal{N}} Z) \\ &= 1 + \sum_i \text{tr } Z_i \mathcal{N} (|i\rangle\langle i|), \end{aligned} \tag{42}$$

where we have made use of the identity $\text{tr}_B ((X \otimes \mathbb{1}) J_{\mathcal{N}}) = \mathcal{N} (X^T)$, and the fact that \mathcal{N} is CPTP in the second line of Eq. (42). As X no longer appears in the objective function, and dephasing it does not affect the constraints, we may assume $X = \sum_i p_i |i\rangle\langle i|$, where p_i are probabilities, without loss of generality. Therefore, we can write

$$\begin{aligned} &\max \left\{ \sum_i p_i + \sum_i \text{tr } Z_i \mathcal{N} (|i\rangle\langle i|) : p_i \mathbb{1} + Z_i \geq 0 \forall i \right\} \\ &= \max \sum_i p_i \text{tr } S_i \mathcal{N} (|i\rangle\langle i|) \\ &= \max \sum_i p_i (1 + C_R(\mathcal{N}(|i\rangle\langle i|))) \\ &= \max_i \{1 + C_R(\mathcal{N}(|i\rangle\langle i|))\}, \end{aligned} \tag{43}$$

where $S_i = \mathbb{1} + p_i^{-1} Z_i \geq 0$ and $\langle j|S_i|j \rangle = 1$ for all j , and we have used the dual form of robustness for states, Eq. (31). Moreover, since the latter is multiplicative under tensor product, see Lemma 18 in Sec. A.1, also the robustness of channels is.

The logged version of the robustness of coherence follows directly from Eq. (43), and the equality of the log-robustness of coherence to the cohering power, $P_R(\mathcal{N})$, of Eq. (5) is also evident from the definition of the latter. To see that the log-robustness of coherence for a channel is also equal to $\hat{P}(\mathcal{N})$, observe that for any $\rho \otimes \omega$

$$\begin{aligned} C_{LR}(\rho \otimes \omega) &\geq C_{LR}(\mathcal{M}(\rho \otimes \omega)) = C_{LR}(\mathcal{N}(\rho) \otimes \sigma) \\ C_{LR}(\rho) + C_{LR}(\omega) &\geq C_{LR}(\mathcal{N}(\rho)) + C_{LR}(\sigma) \\ C_{LR}(\omega) - C_{LR}(\sigma) &\geq C_{LR}(\mathcal{N}(\rho)) - C_{LR}(\rho). \end{aligned} \quad (44)$$

Thanks to Corollary 7, the minimization of the right-hand side of Eq. (44) over ω and σ yields the amortized cost (Eq. (9)), which is equal to the log-robustness of coherence of \mathcal{N} via Theorem 8. Moreover, the inequality in Eq. (44) holds for all $\rho \in \mathcal{H}$. Thus

$$C_{LR}(\mathcal{N}) \geq \max_{\rho \in \mathcal{S}(\mathcal{H})} (C_{LR}(\mathcal{N}(\rho)) - C_{LR}(\rho)) \geq P_R(\mathcal{N}), \quad (45)$$

where the last inequality holds because the maximization in $\hat{P}_R(\mathcal{N})$ is over a larger convex set than that of $P_R(\mathcal{N})$. As the upper and lower bounds on the coherence power are equal it follows that $C_{LR}(\mathcal{N}) = \hat{P}_R(\mathcal{N})$. This completes the proof. \blacksquare

Similarly, one can write the smoothed robustness of coherence (Definition 1) as an SDP in primal and dual form. Using the dual-SDP formulation of the diamond norm [64], the primal SDP of Eq. (2) reads

$$\begin{aligned} 1 + C_R^\epsilon(\mathcal{N}) &= \min \lambda \\ \text{s.t. } J_{\mathcal{M}} &\geq J_{\mathcal{L}}, \\ \text{tr}_B J_{\mathcal{M}} &= \lambda \mathbb{1}_A, \\ \text{tr } J_{\mathcal{M}}(|ij\rangle\langle ik|) &= 0, \quad \forall i \forall j \neq k, \\ V &\geq J_{\mathcal{L}} - J_{\mathcal{N}}, \\ \text{tr}_B V &\leq \epsilon \mathbb{1}_A, \\ \text{tr}_B J_L &= \mathbb{1}_A, \\ J_L &\geq 0, V \geq 0. \end{aligned} \quad (46)$$

The first three constraints correspond to the simulation of channel \mathcal{L} by MIO and the fourth and fifth constraints capture the diamond norm constraints that \mathcal{L} should be ϵ close to \mathcal{N} . The dual form of Eq. (46) is given by

$$\begin{aligned} \max \text{tr} &(J_{\mathcal{N}}(S - \Delta - \epsilon \Omega \otimes \mathbb{1})) \\ \text{s.t. } S &\equiv Y \otimes \mathbb{1} - \sum_i |i\rangle\langle i| \otimes Z_i \geq 0 \\ S - \Delta &\leq W \otimes \mathbb{1} \leq S - \Delta + \Omega \otimes \mathbb{1} \\ \text{tr } Z_i &= 0, \quad \forall i \\ \text{tr } W &= 0, \text{tr } Y = 1 \\ \Delta &\geq 0, \Omega \geq 0, Y \geq 0. \end{aligned} \quad (47)$$

Finally let us state the connection between the log-robustness of coherence of a channel (correspondingly its smoothed version) to the maximum relative entropy (correspondingly ϵ -maximum relative entropy) of the latter with respect to a MIO:

Definition 19 *The maximum relative entropy of two channels \mathcal{N} and \mathcal{M} is:*

$$\begin{aligned} D_{\max}(\mathcal{N}||\mathcal{M}) &:= -\log \max p \\ &\quad \text{s.t., } \mathcal{M} = p\mathcal{N} + (1-p)\mathcal{N}' \\ &\quad p \in [0, 1], \mathcal{N}' \in \text{CPTP}, \\ &= \log \min \{ \lambda : \mathcal{N} \leq \lambda\mathcal{M} \}. \end{aligned} \quad (48)$$

The smoothed version of this quantity is:

$$D_{\max}^{\epsilon}(\mathcal{N}||\mathcal{M}) := \min \left\{ D_{\max}(\mathcal{L}||\mathcal{M}) : \frac{1}{2}\|\mathcal{N} - \mathcal{L}\|_{\diamond} \leq \epsilon \right\}. \quad (49)$$

Proposition 20 *The ϵ -log-robustness of a channel is given by*

$$C_{LR}^{\epsilon}(\mathcal{N}) = \min \{ D_{\max}^{\epsilon}(\mathcal{N}||\mathcal{M}) : \mathcal{M} \text{ MIO} \}. \quad (50)$$

A.3 Channel implementation with maximally coherent resources

In this subsection we prove the main results related to channel implementation with cosdits, Sec. 1.1. We begin with the implementation cost without recycling:

Proof of Theorem 4. Let $\mathcal{L} := \mathcal{M}(\Psi_k \otimes \cdot) : A \rightarrow B$ be the MIO simulation of the induced channel $\mathcal{L} : A \rightarrow B$ such that $\frac{1}{2}\|\mathcal{N} - \mathcal{L}\|_{\diamond} \leq \epsilon$. Noting that $\Psi_k + (k-1)\sigma = \mathbf{1}$, with $\sigma = \frac{\mathbf{1} - \Psi_k}{k-1}$, define

$$\begin{aligned} \mathcal{M}' &:= \mathcal{M}\left(\frac{\mathbf{1}}{k} \otimes \cdot\right) = \frac{1}{k}\mathcal{M}((\Psi_k + (k-1)\sigma) \otimes \cdot) \\ &= \frac{1}{k}\mathcal{M}(\Psi_k \otimes \cdot) + \left(1 - \frac{1}{k}\right)\mathcal{M}(\sigma \otimes \cdot) \\ &= \frac{1}{k}\mathcal{L} + \left(1 - \frac{1}{k}\right)\mathcal{L}', \end{aligned} \quad (51)$$

with $\mathcal{L}' = \mathcal{M}(\sigma \otimes \cdot)$. As \mathcal{M} is MIO, so is $\mathcal{M}' : A \rightarrow B$, and the right-hand-side of Eq. (51) corresponds to a convex decomposition of \mathcal{M}' in terms of $\mathcal{L}, \mathcal{L}' \in \text{CPTP}$. Hence, from the definition of the maximum-relative entropy, Eq. (48) it follows

$$\log k \geq D_{\max}^{\epsilon}(\mathcal{L}||\mathcal{M}') \geq C_{LR}^{\epsilon}(\mathcal{N}). \quad (52)$$

For the converse, let $k \geq 1 + C_R^{\epsilon}(\mathcal{N})$ be an integer. By Eq. (50), it follows that there exists some CPTP map \mathcal{L} with $\frac{1}{2}\|\mathcal{N} - \mathcal{L}\|_{\diamond} \leq \epsilon$ and another CPTP map \mathcal{L}' , such that $(\mathcal{L} + (k-1)\mathcal{L}')/k$ is MIO. Make the following ansatz for a channel \mathcal{M} that is feasible for the simulation cost:

$$\mathcal{M}(\tau \otimes \rho) := \text{tr}(\Psi_k \tau) \mathcal{L}(\rho) + \text{tr}((\mathbf{1} - \Psi_k) \tau) \mathcal{L}'(\rho).$$

The map \mathcal{M} is MIO if and only if $\mathcal{M}(|i\rangle\langle i| \otimes \cdot)$ is MIO for all incoherent basis states $|i\rangle$. This is the case, since $\mathcal{M}(|i\rangle\langle i| \otimes \cdot) = \frac{1}{k}\mathcal{L} + (1 - \frac{1}{k})\mathcal{L}'$, which is MIO by construction. Hence $\log k \geq C_{\text{sim}}^{\epsilon}(\mathcal{N})$ and the former can be taken as small as $\lceil 1 + C_R^{\epsilon}(\mathcal{N}) \rceil$. As the implementation cost is necessarily an integer Eq. (7) follows. This completes the proof. \blacksquare

The proof for the simulation cost with recycling relies on the previous result, effectively implementing a simulation of the target channel and the resource.

Proof of Theorem 6. Let $\mathcal{T} : \mathbb{C} \otimes A \rightarrow S \otimes B$ such that $\mathcal{T}(1 \otimes \rho) = \sigma \otimes \mathcal{N}(\rho)$. Here, we have made use of the fact that any state can be identified with a preparation channel from \mathbb{C} to S mapping the unique state $1 \rightarrow \sigma$. Now let \mathcal{M} be a MIO such that

$$\mathcal{M}(\Psi_k \otimes \cdot) = \mathcal{T}(\cdot).$$

From Theorem 4 such a simulation is possible if and only if

$$\begin{aligned} k &\geq (1 + C_R(\mathcal{T})) \\ &= (1 + C_R(\sigma))(1 + C_R(\mathcal{N})) \end{aligned}$$

where we have used the multiplicativity of the robustness of coherence in the last line. Taking the logarithm on both sides and re-arranging terms gives Eq. (10). This completes the proof. \blacksquare

We will now make use of Theorem 6 to prove Theorem 8 regarding the amortized cost.

Proof of Theorem 8. From Definition 5 and Corollary 7 it follows immediately that the zero-error amortized cost is equal to the log-robustness. Indeed, if we take the log on each side of Eq. (11) and then let $k, m \rightarrow \infty$ we obtain $C_{\text{amo}}(\mathcal{N}) = C_{LR}(\mathcal{N})$. This in turn implies that the ϵ -error amortized cost, Eq. (9), can be rewritten in terms of the log-robustness of the channel \mathcal{L} as

$$C_{\text{amo}}^\epsilon(\mathcal{N}) = \min C_{LR}(\mathcal{L}) \quad \text{s.t.} \quad \frac{1}{2} \|\mathcal{N} - \mathcal{L}\|_\diamond \leq \epsilon \equiv C_{LR}(\mathcal{N}). \quad (53)$$

\blacksquare

A.4 Channel implementation with arbitrary resources: SDPs

Here we prove the SDP forms of the generic implementation problem considered in Sec. 1.2. We begin by first providing an SDP for the optimal implementation of unitary gates without recycling using the gate fidelity as a figure of merit. We then proceed to prove Proposition 10. For ease of notation, we swap the output subspaces of the MIO implementation map throughout this section, i.e., $\mathcal{M} : R \otimes A \rightarrow B \otimes S$. We also consider unnormalized Choi matrices.

Proposition 21 *The optimal gate fidelity of implementation of a unitary $U : A \rightarrow B$ by means of a MIO $\mathcal{M} : R \otimes A \rightarrow B \otimes S$ and a pure coherent state $\omega \in R$ is given by the following SDP:*

$$\begin{aligned} F &= \max \frac{1}{d_A^2} \text{tr} \left((\omega^T \otimes J_U \otimes \mathbf{1}_S) X \right) \\ &\quad \text{s.t. } X \text{ is the Choi matrix of a MIO,} \end{aligned} \quad (54)$$

where J_U is the Choi matrix of the channel $U \cdot U^\dagger$, $d_A = \dim(A)$, and ω^T denotes the transpose of ω .

Proof. Consider a maximally incoherent operation $\mathcal{M} : R \otimes A \rightarrow B \otimes S$. Its Choi matrix is defined as

$$\begin{aligned} J_{\mathcal{M}} &= (\text{id}_{RA} \otimes \mathcal{M})(\Phi_{RR'} \otimes \Phi_{AA'}) \\ &= \sum_{i,k,j,l} |ik\rangle_{RA} \langle jl| \otimes \mathcal{M}(|ik\rangle_{R'A'} \langle jl|). \end{aligned} \quad (55)$$

Suppose that $\mathcal{E} = \text{tr}_S(\mathcal{M}(\omega \otimes \cdot))$, where $\mathcal{E} : A \rightarrow B$. Then the Choi matrix of \mathcal{E} , $J_{\mathcal{E}}$, is explicitly given by

$$\begin{aligned} J_{\mathcal{E}} &= \text{tr}_{RS} \left((\omega^T \otimes \mathbf{1}_A \otimes \mathbf{1}_B \otimes \mathbf{1}_S) J_{\mathcal{M}} \right) \\ &= \sum_{ijkl} \langle j|\omega^T|i\rangle \langle k\rangle_A \langle l| \otimes \text{tr}_S(\mathcal{M}(|ik\rangle_{R'A'} \langle jl|)) \\ &= \sum_{k,l} |k\rangle_A \langle l| \otimes \text{tr}_S(\mathcal{M}(\omega \otimes |k\rangle_{A'} \langle l|)) \\ &= (\text{id}_A \otimes \mathcal{E})(\Phi_{AA'}) \end{aligned} \quad (56)$$

On the other hand, the Choi matrix of the unitary channel is given by $J_U = (\mathbb{1}_A \otimes U) \Phi_{AA'} (\mathbb{1}_A \otimes U^\dagger)$.

Our aim is to compute how well the map \mathcal{E} implements the unitary channel $\mathcal{N}(\rho) = U\rho U^\dagger$. To that end, we use the *gate fidelity*, i.e., the fidelity between the Choi matrices of the corresponding channels:

$$\begin{aligned} F &= \frac{1}{d_A^2} \mathbf{tr}(J_U J_{\mathcal{E}}) \\ &= \frac{1}{d_A^2} \mathbf{tr} \left((\mathbb{1}_R \otimes J_U \otimes \mathbb{1}_S) \left(\omega^T \otimes \mathbb{1}_A \otimes \mathbb{1}_B \otimes \mathbb{1}_S \right) J_{\mathcal{M}} \right) \\ &= \frac{1}{d_A^2} \mathbf{tr} \left(\left(\omega^T \otimes J_U \otimes \mathbb{1}_S \right) J_{\mathcal{M}} \right). \end{aligned} \quad (57)$$

In particular, we want to obtain the optimal gate fidelity of implementation of the given channel. This is an SDP in primal standard form, where $J_{\mathcal{M}} \geq 0$ is the semidefinite variable subject to the constraints $\mathbf{tr}_{BS} J_{\mathcal{M}} = \mathbb{1}_R \otimes \mathbb{1}_A$ and $\mathbf{tr}((|ik\rangle_{RA}\langle ik| \otimes |jl\rangle_{SB}\langle j'l'|) J_{\mathcal{M}}) = 0$ for all i, k and all $j \neq j', l \neq l'$. ■

In the case of a generic channel \mathcal{N} , the diamond norm distance needs to be used, as discussed in the main text.

Proof of Proposition 10. The SDP can now be formulated as follows:

$$\begin{aligned} \min \quad & \frac{1}{2} \|\mathcal{E} - \mathcal{N}\|_{\diamond} \\ \text{s.t.} \quad & X \text{ is the Choi matrix of a MIO } \mathcal{M}, \\ & \mathcal{E} = \mathbf{tr}_B \mathcal{M}(\omega \otimes \cdot). \end{aligned} \quad (58)$$

Using the result of [64], the dual form of the diamond norm distance can be written as

$$\begin{aligned} \min \quad & \lambda \\ \text{s.t.} \quad & Z \geq J_{\mathcal{E}} - J_{\mathcal{N}} \\ & \mathbf{tr}_B Z \leq \lambda \mathbb{1}_A \\ & Z \geq 0. \end{aligned} \quad (59)$$

Using Eq. (56) to relate the Choi matrix of \mathcal{E} to that of \mathcal{M} is MIO it is then straightforward to obtain the final formulation. ■

A.5 Alternative definitions of coherence left

In this appendix we discuss possible bounds on and alternative definitions of coherence left, when an arbitrary resource is used to simulate a channel with recycling. Let us start by relaxing the tensor-product constraint at the output of the implementation: a simpler problem is that of finding a MIO that simulates the whole tensor-product output up to a given error. This also constitutes an upper bound on the robustness of coherence left.

Proposition 22 *The maximum robustness of coherence left in the resource $\sigma \in S$ after the implementation of a quantum channel $\mathcal{N} : A \rightarrow B$ via MIO $\mathcal{M} : R \otimes A \rightarrow S \otimes B$ and a coherent resource $\omega \in R$ up to error ϵ in diamond norm, i.e., Eq. (20), is bounded from above by the following optimization problem:*

$$\begin{aligned} \max \quad & C_R(\sigma) \\ \text{s.t.} \quad & \frac{1}{2} \|\sigma \otimes \mathcal{N} - \mathcal{M}(\omega \otimes \cdot)\|_{\diamond} \leq \epsilon, \end{aligned} \quad (60)$$

which can be expressed as the following semidefinite program:

$$\begin{aligned}
& \max \sum_{i \neq j} \sigma_{ij} \\
& \text{s.t. } J_{\mathcal{M}} \text{ is the Choi matrix of } \mathcal{M} \text{ MIO} \\
& Z \geq \sigma \otimes J_{\mathcal{N}} - \mathbf{tr}_R \left(\left(\omega^T \otimes \mathbf{1}_{SAB} \right) J_{\mathcal{M}} \right) \\
& \mathbf{tr}_{SB} Z \leq \epsilon \mathbb{1}_A \\
& Z \geq 0.
\end{aligned} \tag{61}$$

where $J_{\mathcal{N}}$ is the Choi matrix of \mathcal{N} .

Proof. It is straightforward to show that Eq. (60) is an upper bound on Eq. (20). Indeed, we can make the separable-output ansatz $\mathcal{M}(\omega \otimes \cdot) = \sigma \otimes \mathcal{L}(\cdot)$ for the MIO \mathcal{M} simulating the channel in Eq. (60) and obtain the problem of Eq. (20).

As for the SDP formulation of this problem, the last three conditions in Eq. (61) are a simple translation of the diamond norm error one in Eq. (60), following [64]. It then remains to show that the robustness of coherence can be re-cast as the sum of off-diagonal elements of $\sigma \in S$. To that end recall the dual formulation, Eq. (31), of the robustness of coherence which reads

$$1 + C_R(\sigma) = \max \mathbf{tr} \sigma S = \max \mathbf{tr} \left(\sigma \circ S^T \right) G, \tag{62}$$

where S is a non-negative-definite matrix with ones on the diagonal and we have introduced the all-one matrix $G_{ij} = 1$ for all i, j and the Hadamard component-wise product $\sigma \circ S = \sum_{ij} \sigma_{ij} S_{ij} |i\rangle\langle j|$. That this is equal to the sum of all off-diagonal elements of σ follows from

$$\begin{aligned}
C_R(\sigma) &= \sum_{i,j} \sigma_{ij} S_{ji} - 1 \\
&= \sum_i \sigma_{ii} + \sum_{i \neq j} \sigma_{ij} S_{ji} - 1 \\
&= \sum_{i \neq j} \sigma_{ij} S_{ji},
\end{aligned} \tag{63}$$

where we have made use of the fact that $S_{ii} = 1 \forall i$ in the second line above. As the transformation $\sigma \rightarrow \sigma \circ S^T$ can be implemented by a MIO CPTP map, given that $(\sigma \circ S^T)_{ij} \leq \sigma_{ij} \forall i \neq j$ (no coherence is created in the transformation), the two objective functions are equivalent under the given constraints, since the diamond norm is contractive under CPTP maps. ■

The SDP of Proposition 22 amounts to requiring that the MIO \mathcal{M} implements a good simulation of the overall output $\sigma \otimes \mathcal{N}$ when provided with the input resource ω . Note that this implies that the local reduced output systems of the implementation \mathcal{M} are close to σ and \mathcal{N} , but also that their correlations are small, so that the implementation is ϵ -close to a tensor-product one. However, if we increase the allowed dimension of σ it is possible to find states that are ϵ -close to it but have an increasing amount of coherence, e.g., $(1 - \epsilon) \sigma + \epsilon \Psi_d$. Hence Eq. (60) can give an unbounded amount of coherence left if the output dimension of the resource is allowed to increase.

Two possible ways to remedy this involve changing the objective function such that it either maximizes the ϵ -robustness of coherence of the output resource state, or to find the most coherent state among all output resource states that are ϵ -close to the desired target. Both of these solutions provide an alternative definition of coherence left, physically motivated by the fact that the true output resource is not σ but only a state ϵ -close to it, either in robustness or in trace norm. Unfortunately, the former is not an SDP while the latter is not easily computable. Which of these, or other, definitions of coherence left is better may depend on the recycler's objective and we leave it as an open question for the interested reader.

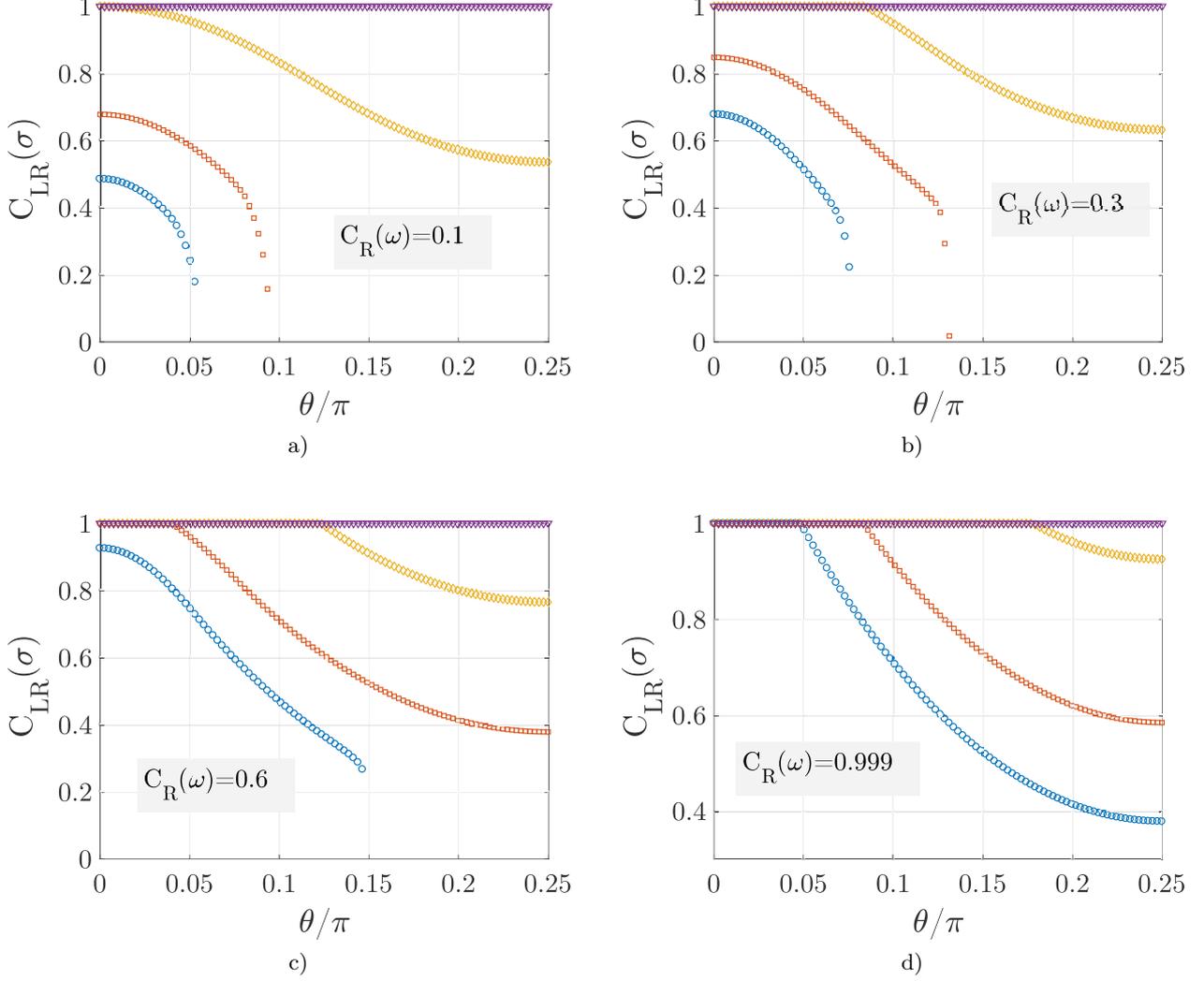


Figure 6: Maximum robustness of coherence left in the resource after the approximate MIO implementation of the qubit unitary U_θ vs. θ/π for four values of input coherence (a-d). Each plot shows four curves corresponding to different error thresholds $\epsilon \in \{0.15, 0.25, 0.45, 0.75\}$ (from lower to higher curves).

In Fig. 6 we plot the robustness of coherence left after the approximate implementation of a qubit unitary U_θ vs. θ/π , for several values of the input coherence and of the error threshold. As expected, for a fixed input coherence, more coherence can be recovered at the output if we accept a worse implementation. Moreover, for sufficiently small error thresholds, there are unitaries that cannot be implemented at all.

A.6 State conversion under MIO and flagpole states

In this subsection we show that any pure resource state in dimension larger than 2 can implement some coherent unitary channel. We will first establish a necessary (but not sufficient) criterion for state conversion under MIO (Lemma 12), and then show that there exist families of flagpole states (see Definition 13), not convertible to cosdits, that can be used as resources for implementing *any* quantum channel. In particular, we will demonstrate that for any channel $\mathcal{N} : A \rightarrow B$, there exists $\mathcal{M} : R \otimes A \rightarrow B$, \mathcal{M} MIO and a particular family of pure flagpole states $|\phi_p\rangle \in R$ capable of simulating the channel (Theorem 14).

Proof of Lemma 12. The geometric measure of coherence, introduced in [60], is defined as

$$C_g(\rho) = 1 - F_C(\rho) \quad (64)$$

where

$$F_C(\rho) = \max_{\delta \in \Delta} F(\rho, \delta)^2, \quad (65)$$

is the familiar Uhlmann fidelity between two mixed states, $F(\rho, \delta) = \|\sqrt{\rho}\sqrt{\delta}\|_1$. In particular for a pure state $\rho = |\psi\rangle\langle\psi|$, $F(\psi, \delta)^2 = \text{tr} \psi\delta$, and $F_C(\psi) = \lambda_1(\psi) = \max_i \langle i|\psi|i\rangle$. As the geometric measure of coherence is monotonically decreasing under MIO, it follows that $F_C(\mathcal{M}(\rho)) \geq F_C(\rho)$, for all \mathcal{M} MIO. Setting $\phi = \mathcal{M}(\psi)$ leads to the desired result. ■

Proof of Theorem 14. Let us decompose the state space R into the following orthogonal subspaces

$$\varphi_p, \phi, \Pi = \mathbb{1} - \varphi_p - \phi \quad (66)$$

where $\varphi_p = |\varphi_p\rangle\langle\varphi_p|$, with $|\varphi_p\rangle$ given by Eq. (22), ϕ_1 is the rank-one projector onto the state

$$|\phi\rangle = \sqrt{1-p}|0\rangle - \sqrt{\frac{p}{d-1}} \sum_{j=1}^{d-1} |j\rangle, \quad (67)$$

and Π is a rank- $(d-2)$ orthogonal complement to φ_p and ϕ .

We now make the following ansatz for the MIO channel simulating \mathcal{N} with a flagpole resource φ_p :

$$\begin{aligned} \mathcal{M}(\rho \otimes \sigma) = & \mathcal{N}(\rho) \text{tr}(\sigma\varphi_p) + \mathcal{L}_1(\rho) \text{tr}(\sigma\phi_1) \\ & + \mathcal{L}_2(\rho) \text{tr}(\sigma\Pi), \end{aligned} \quad (68)$$

where $\mathcal{L}_{1,2}$ are CPTP maps. In order for \mathcal{M} to be MIO we require that $\mathcal{M}(\cdot \otimes |j\rangle\langle j|)$ is MIO, $\forall j \in (0, \dots, d-1)$. This leads to the following two conditions

$$\begin{aligned} p\mathcal{N} + (1-p)\mathcal{L}_1 \text{ is MIO} & \quad \text{for } j = 0 \\ \frac{1-p}{d-1}\mathcal{N} + \frac{p}{d-1}\mathcal{L}_1 + \left(1 - \frac{1}{d-1}\right)\mathcal{L}_2 \text{ is MIO} & \quad \text{for } j > 0 \end{aligned} \quad (69)$$

From Proposition 19 we know that there exists a CPTP map \mathcal{L}_1 such that the first condition in Eq. (69) is fulfilled if and only if $p \leq (1 + C_R(\mathcal{N}))^{-1}$. To prove that the second condition in Eq. (69) is satisfied note that it may be written as

$$\frac{1}{d-1}\mathcal{N}' + \left(1 - \frac{1}{d-1}\right)\mathcal{L}_2 \quad (70)$$

where we defined the CPTP map $\mathcal{N}' = (1-p)\mathcal{N} + p\mathcal{L}_1$ whose robustness of coherence is at most $d-1$. By Definition 19 this implies that there exists a CPTP map \mathcal{L}_2 such that $(\mathcal{N}' + (d-2)\mathcal{L}_2)/(d-1)$ is MIO. This completes the proof. ■

A.7 Qubit unitaries

In this appendix we prove Propositions 15 and 16 pertaining to the implementation of unitary gates via MIO using non-maximally coherent qubit states and d -dimensional flagpole states respectively. We first prove that exact implementation of unitary gates via MIO can only be achieved using cosbit resources, and then establish the subset of d -dimensional flagpole states that allow for the implementation of any qubit unitary via MIO.

Proof of Proposition 15. Let $\{|\omega\rangle, |\omega^\perp\rangle\}$ be an orthonormal basis for \mathbb{C}^2 , where

$$\begin{aligned} |\omega\rangle &= \alpha|0\rangle + \beta|1\rangle \\ |\omega^\perp\rangle &= \beta|0\rangle - \alpha|1\rangle \end{aligned} \quad (71)$$

with $\alpha, \beta \in \mathbb{C}$, $|\alpha|^2 + |\beta|^2 = 1$. The Kraus decomposition of the most general \mathcal{M} is MIO implementing any unitary operator $U : \mathcal{H}_d \rightarrow \mathcal{H}_d$, is given by

$$K_i = \lambda_i U \otimes |0\rangle\langle\omega| + R^i \otimes |0\rangle\langle\omega^\perp|, \quad (72)$$

where $\{\lambda_i \in \mathbb{C}\}$, and $\{R^i : \mathcal{H}_d \rightarrow \mathcal{H}_d, 1 \leq d' \leq d\}$. As \mathcal{M} is CPTP, $\sum_i K_i^\dagger K_i = \mathbb{1}$ which yields the following conditions on $\{\lambda_i, R^i\}$

$$\begin{aligned} \sum_i |\lambda_i|^2 &= 1 \\ \sum_i |\lambda_i|^2 R^i &= 0 \\ \sum_i |\lambda_i|^2 R^{i\dagger} R^i &= \mathbb{1}. \end{aligned} \quad (73)$$

Applying them together with the MIO condition

$$\langle m | \left(\sum_i K_i |jk\rangle \langle jk| K_i^\dagger \right) |n\rangle = 0, \quad \forall m \neq n$$

gives rise to the following two equations

$$|\alpha|^2 U_{mj} U_{jn}^\dagger + |\beta|^2 \sum_i |\lambda_i|^2 R_{mj}^i (R^{i\dagger})_{jn} = 0 \quad (74)$$

$$|\beta|^2 U_{mj} U_{jn}^\dagger + |\alpha|^2 \sum_i |\lambda_i|^2 R_{mj}^i (R^{i\dagger})_{jn} = 0, \quad (75)$$

corresponding to $k = 0$ and $k = 1$ respectively. Summing both equations above we obtain the additional condition $\sum_i |\lambda_i|^2 R_{mj}^i (R^{i\dagger})_{jn} = -U_{mj} U_{jn}^\dagger$ which, upon substituting into any one of Eqs. (74, 75) results in

$$\left(|\alpha|^2 - |\beta|^2 \right) U_{mj} U_{jn}^\dagger = 0, \quad \forall j, n, m. \quad (76)$$

Now, unless U is the identity, there is at least one pair of values (j, n) such that $0 < |U_{nj}| < 1$. Moreover, as $\sum_m |U_{mj}|^2 = 1$ there also exists at least one $m \neq n$ such that $|U_{mj}| > 0$. Thus for Eq. (76) to hold true for all j, n, m , $|\alpha|^2 = |\beta|^2 = \frac{1}{2}$ which implies that $|\omega\rangle = |\Psi_2\rangle$. This completes the proof. \blacksquare

Proof of Proposition 16. We shall first prove the statement of the Proposition for the case of qutrit flagpoles. To that end consider the following orthonormal basis for \mathbb{C}^3 ,

$$\begin{aligned} |\varphi_p\rangle &\equiv |\phi_0\rangle = \sqrt{p}|0\rangle + \sqrt{\frac{1-p}{2}}(|1\rangle + |2\rangle), \\ |\phi_1\rangle &= \frac{1}{\sqrt{2}}(|1\rangle - |2\rangle), \\ |\phi_2\rangle &= \sqrt{1-p}|0\rangle - \sqrt{\frac{p}{2}}(|1\rangle + |2\rangle). \end{aligned}$$

As in the previous case, the Kraus decomposition of the most general MIO \mathcal{M} implementing $U_\theta : \mathbb{C}^2 \rightarrow \mathbb{C}^2$ is given by

$$K_i = \lambda_i U_\theta \otimes |0\rangle\langle\phi_0| + R^i \otimes |0\rangle\langle\phi_1| + P^i \otimes |0\rangle\langle\phi_2|. \quad (77)$$

As \mathcal{M} is CPTP, $\sum_i K_i^\dagger K_i = \mathbb{1}$, which imposes the following conditions on $\{\lambda_i, R^i, P^i\}$

$$\sum_i |\lambda_i|^2 = 1 \quad (78)$$

$$\sum_i |\lambda_i|^2 R^i = \sum_\alpha P_\alpha = 0 \quad (79)$$

$$\sum_i |\lambda_i|^2 R^{i\dagger} R^i = \sum_i |\lambda_i|^2 P^{i\dagger} P^i = \mathbb{1} \quad (80)$$

$$\sum_i |\lambda_i|^2 R^{i\dagger} P^i = \sum_i |\lambda_i|^2 P^{i\dagger} R^i = 0. \quad (81)$$

Applying the above conditions together with the MIO condition

$$\langle m | \left(\sum_{\alpha} K_{\alpha} |jk\rangle \langle jk| K_{\alpha}^{\dagger} \right) |n\rangle = 0$$

for $m \neq n = 0, 1$, $j = 0, 1$ and $k = 0, 1, 2$ we obtain

$$U_{mj} U_{jn}^{\dagger} |\phi_0^k|^2 + A_j |\phi_1^k|^2 + B_j |\phi_2^k|^2 + C_j \phi_1^k \phi_2^k = 0, \quad (82)$$

where we have defined

$$\begin{aligned} A_j &= \sum_i |\lambda_i|^2 R_{mj}^i \left(R^{i\dagger} \right)_{jn} \\ B_j &= \sum_i |\lambda_i|^2 P_{mj}^i \left(P^{i\dagger} \right)_{jn} \\ C_j &= \sum_i |\lambda_i|^2 \left(R_{mj}^i \left(P^{i\dagger} \right)_{jn} + P_{mj}^i \left(R^{i\dagger} \right)_{jn} \right), \end{aligned}$$

and $\phi_l^k \equiv \langle \phi_l | k \rangle$. Fixing $m = 0$, $n = 1$, Eq. (82) gives rise to six equations which we can solve for A_j , B_j and C_j to obtain

$$\begin{aligned} A_0 &= -A_1 = -\frac{(2p-1)}{p-1} \cos \theta \sin \theta \\ B_0 &= -B_1 = \frac{p}{p-1} \cos \theta \sin \theta \\ C_0 &= C_1 = 0. \end{aligned} \quad (83)$$

Now by definition A_j and B_j correspond to the off-diagonal elements of CPTP maps acting on the state $|j\rangle\langle j|$. As the resulting operator must be positive it follows that $|A_j| \leq \frac{1}{2}$, $|B_j| \leq \frac{1}{2}$. Moreover, the conditions on the A_j in Eq. (83) are satisfied if B_0, B_1 are given as in Eq. (83). Hence

$$|B_j| = \frac{p}{1-p} \cos \theta \sin \theta \leq \frac{1}{2},$$

which implies

$$p \leq \frac{1}{1 + \sin 2\theta}.$$

Therefore, there can be no MIO that implements the qubit unitary, U_{θ} , with a qutrit flagpole state $|\varphi_p\rangle$ such that $p > (1 + \sin 2\theta)^{-1}$.

The proof can be extended to arbitrary d -dimensional flagpole and $(d-1)$ -dimensional unitary. We note, however, that in this case we obtain the following upper bound $p \leq \left(1 + \frac{C_R(U)}{d-2}\right)^{-1}$, which coincides with Eq. (24) only for $d = 3$. This completes the proof. \blacksquare

We note that it is an interesting open question whether less coherent flagpole states allow for the implementation of qubit unitaries via MIO for the case $d > 3$.

A.8 Asymptotic MIO generation capacity

In this appendix we prove Thm. 17, namely that the asymptotic coherence-generating capacity of a channel under MIOs is equal to its complete relative-entropy cohering power. The proof is a straightforward extension of [5, Thm. 1], where IOs were considered instead. The key difference here is that the coherence cost of arbitrary states under MIOs is given by the relative entropy of coherence [67], whereas under IOs this is true only for pure states.

Proof of Theorem 17. We first prove the upper bound: for any coherence-generating protocol, the trace-distance between the final state ρ_n and the cosdit can be upper bounded as $\|\rho_n - \Psi_2^{\otimes nR}\|_1 \leq 2\sqrt{\epsilon}$. Then the asymptotic continuity of the relative-entropy of coherence [65, Lemma 12] implies that:

$$\left| C_r \left(\Psi_2^{\otimes nR} \right) - C_r (\rho_n) \right| \leq 2\sqrt{\epsilon} nR + 2h(\sqrt{\epsilon}), \quad (84)$$

where $h(x) = -x \log x - (1-x) \log(1-x)$ is the binary-entropy function. Hence we obtain the following chain of inequalities:

$$\begin{aligned} nR(1 - 2\sqrt{\epsilon}) - 2h(\sqrt{\epsilon}) &\leq C_r(\rho_n) \\ &= \sum_{j=0}^{n-1} (C_r(\rho_{j+1}) - C_r(\rho_j)) \\ &\leq n \sup_{\rho \in AC} (C_r((\mathcal{N} \otimes \text{id})(\rho)) - C_r(\rho)), \end{aligned} \quad (85)$$

where the first inequality follows from Eq. (84) and $C_r(\Psi_2^{\otimes nR}) = nR$, while the equality by adding and subtracting the relative-entropy of coherence of each intermediate state of the protocol. The last inequality instead follows from substituting each term of the sum with its sup over all states and the monotonicity of C_r under MIOs. After dividing both sides by n and taking the $n \rightarrow \infty$ limit we obtain that the rate of any coherence-generating protocol is upper-bounded by $\mathbb{P}_r(\mathcal{N})$ up to an error vanishing with ϵ .

For the lower bound instead we need to exhibit a protocol that asymptotically attains $\mathbb{P}_r(\mathcal{N})$. The protocol works as follows: i) we first apply the channel to an incoherent state $|0\rangle\langle 0|$ in order to produce some coherence, i.e., $\sigma = \mathcal{N}(|0\rangle\langle 0|)$; ii) we produce a sufficient number of copies of σ to distill a certain amount of cosbits, which are then used to produce a target state ρ via coherence dilution; iii) we apply the channel to ρ in order to obtain a more coherent state $\rho' = (\mathcal{N} \otimes \text{id})(\rho)$; iv) we distill cosbits from ρ' ; v) we use the increased amount of cosbits obtained to iterate the processes (iii-iv) k times. The asymptotic coherence-generating capacity and cost of states under MIOs are both equal to the relative-entropy of coherence [67], so that the conversion rates of processes (ii-iv) described above can be written as

$$\begin{aligned} \sigma^{\otimes m} &\mapsto \Psi_2^{\otimes \lfloor m(C_r(\sigma) - \delta) \rfloor}, \\ \Psi_2^{\otimes \lfloor n(C_r(\rho) + \delta) \rfloor} &\mapsto \rho^{\otimes n} \\ \rho'^{\otimes n} &\mapsto \Psi_2^{\otimes \lfloor n(C_r(\rho') - \delta) \rfloor}, \end{aligned} \quad (86)$$

where $m \gtrsim n(C_r(\rho) + \delta) / (C_r(\sigma) - \delta)$. Note that the coherence production of processes (iii-iv) is not larger than $n(R - 2\delta)$, where $R = C_r(\rho') - C_r(\rho)$, and each transformation is accurate up to an error ϵ . Hence, the overall coherence-production rate is bounded by

$$\frac{kn(R - 2\delta) + n(C_r(\phi) + \delta)}{kn + m} \xrightarrow[k \rightarrow \infty]{} R - 2\delta, \quad (87)$$

which can be made arbitrarily close to R by taking δ sufficiently small. We have therefore described a protocol that generates coherence at an asymptotical rate $C_r((\mathcal{N} \otimes \text{id})(\phi)) - C_r(\phi)$ for any ϕ_{AC} using MIOs. By taking the sup of this quantity, we obtain the desired lower bound. \blacksquare